

Proposal

Nov. 18, 1970

I. OBSERVABLE STATES OF A SUBPROCESS

These are the possible values of the call stack entry for the subprocess, when it is not the current running subprocess. This description does not include F-return count or class code information/

- A. About to execute instruction word at address P.
- B. In the middle of a complicated XJ instruction at address P.
Should occur only on explicit or implicit subprocess call. Implicit subprocess call's occur due to an error return or F-returns with multilevel operations.
- C. Finishing up an XJ instruction at address P.

The body of the action has been completed, a normal return has occurred, but no P counter offset computation has yet occurred, two cases where this can occur are:

- i. When the P-counter offset with place P-counter outside the field length or negative.
- ii. On a get event or luug ECS action when an interrupt is recognized after an event has been received.

II. A CATALOGUE OF VARIOUS CASES

- A. ECS action (other than get event or hang) *luug*

State B cannot occur unless the ECS action returns with F return in a multilevel operation or error returns, (in both cases causes immediate implicit subprocess call.)

- B. Get event or hang

State B can only occur in the event of an error return. If an interrupt is recognized before an event is received, state A occurs iwth P equal to the address of the XJ instruction. If an interrupt is recognized after an event B received then the stack entry will be in state C with P equal to the address of the XJ instruction.

- C. Subprocess call, explicit or implicit (except error due to P-counter offset negative or outside field length)

The stack entry will be in state B with P equal to the address of the XJ instruction. Any interrupts will not be recognized until some other subprocess is actually running.

- D. Interrupts

When an interrupt is recognized the interrupted subprocess entry in the call stack will be set as follows:

- i. If recognized between instruction words, the stack state will become A with P equal to the address of the next instructions to execute.

- ii. If recognized during get event or hang and before an event is received, the stack state will become A with P equal to the address of the get event or hang XJ instruction.
- iii. If recognized during get event or hang and after an event has been received, the stack state will become C with P equal to the address of the get event or hang XJ instruction.
- iv. The above 3 cases are the only cases in which an interrupt can be recognized by a subprocess.

E. Errors

When an error occurs in a subprocess the stack entry will be set as follows:

- i. Arithmetic or scope type error.
The stack entry will be placed in state A
- ii. Error return from ECS action or explicit or implicit subprocess call.
The stack entry will be placed in state B with P equal to the address of the XJ instruction.
- iii. P counter off set would make P counter negative or outside field length.
The stack entry will be placed in state C with P equal to the address of the XJ instruction.
- iv. There is no XJ instruction at the address expected for computation of a P counter offset.
The stack entry will be placed in state C with P equal to the address of the XJ instruction.
- v. The above 4 cases are the only cases in which an error can be recognized. *non-existent*

III RETURN ACTIONS (which depend on the state of the stack entry.)

A. Normal return (included return with parameters)

- i. If the stack was in state A, execution resumes with the instruction at P.
- ii. If the stack was in state B, P counter offset is computed from the XJ instruction at address P, and execution resumes at the computed address.
- iii. If the stack was in state C, same as ii).

B. Special return ("return and re-execute")

- i. If the stack was in state A, execution resumes with the instruction at P.
- ii. If the stack was in state B, same as i)
(This is the re-execute case)
- iii. If the stack was in state C, same as A)ii)above.